Deadlock

Recall: Deadlock

- Situation in which group of threads/processes all block forever
- Typically, each holds a resource that others are blocking on



http://minutillo.com/steve/weblog/2003/1/21/deadlock/, where it is attributed to "Chuck @ China" (http://chake.chinatefl.com/)

Modeling resource contention deadlocks

- Focus on requesting and freeing resources
- Assume process blocks if requests something in use
- Vertices for each process and each resource
- Edge from resource to process holding it, from process to resource it is blocking on

•	Example:	Process A	Process B	Process C
		request R	request S	request T
		request S	request T	request R
		free R	free S	free T
		free S	free T	free R

Detecting deadlock

- Use DFS to look for cycle in graph
- Vertex colors: "unvisited", "in progress", "done"

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all vertices initially "unvisited"
for each "unvisited" vertex v: visit(v);
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```
visit(vertex v):
  color v "in progress";
  for each neighbor u of v:
    if(u is "in progress") print "found cycle";
    if(u is "unvisited") visit(u);
  color v "done";
```

Detection strategy w/ multiple copies of resources

(using vectors and matrices...)

Is the following situation safe?

additional resources needed per job

R =	1	0	2
	2	2	1
	0	1	1
	3	1	0

currently used resources per job

C =	1	1	0
	0	1	1
	2	0	0
	1	0	0

- A. Yes
- B. Yes

- currently available resources
- A = 2 1 1

- C. No
- D. No
- E. As long as you don't call on me to explain my answer

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- A. Yes
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- currently available resources
- A = 2 1 1

- C. <u>No</u>
- D. <u>No</u>
- E. As long as you don't call on me to explain my answer

Strategy 1: Deadlock avoidance ("Banker's Algorithm")

- Eliminate possibility of deadlock through clever resource allocation
- For each resource request, evaluate if granting it puts system into an unsafe state (only grant it if not)

Would the Banker's Algorithm grant a request from job 2 for 1 unit of resource 3 in the following situation?

currently available resources

A = 2 1 1

currently used resources per job				
	1	0	2	
C –	2	2	1	
U –	0	1	1	
	0	1	0	

additional resources needed per job

	1	0	1
R =	3	2	1
η –	0	2	0
	2	1	0

A. Yes

C. None of the above

B. No

Would the Banker's Algorithm grant a request from job 2 for 1 unit of resource 3 in the following situation?

currently available resources

A = 2 1 1

 $C = \begin{bmatrix} 1 & 0 & 2 \\ 2 & 2 & 1 \\ 0 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix}$

additional resources needed per job

R =	1	0	1
	3	2	1
	0	2	0
	2	1	0

C. None of the above

B. No

A. Yes